

The Minimum Rank Problem for Finite Fields

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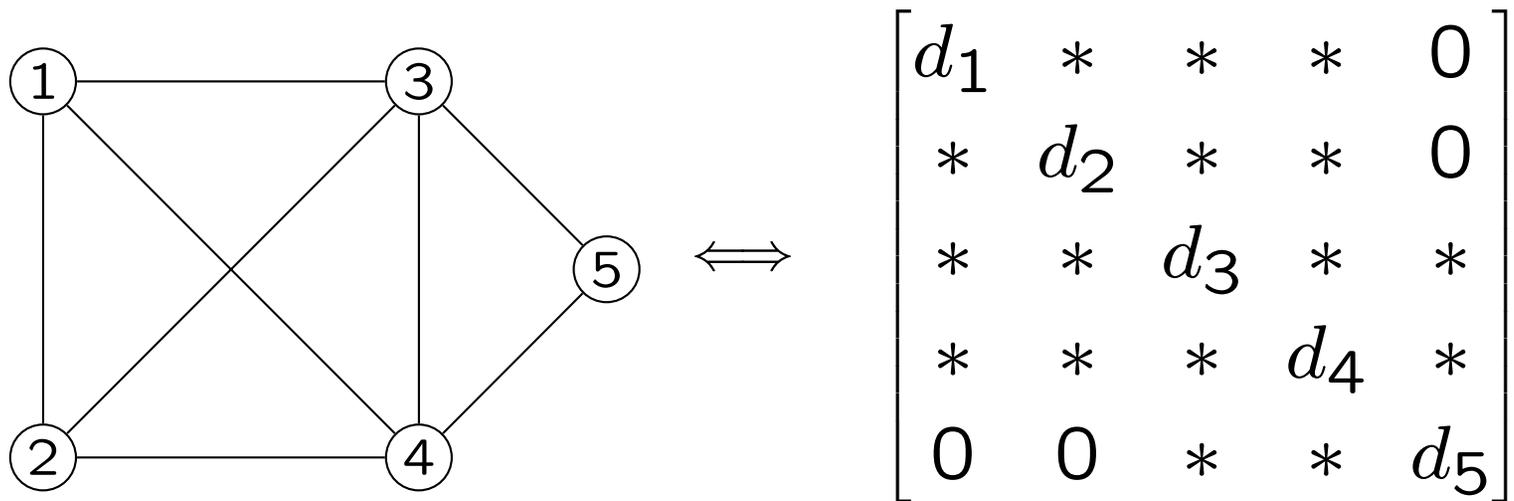
References

Barrett, van der Holst, Loewy, Graphs whose Minimal Rank is Two, *Electronic Journal of Linear Algebra*, volume 11 (2004), pp. 258–280

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Barrett, Grout, March, The Minimal Rank Problem over a Finite Field, in preparation.

Example of $S(F, G)$



$$d_1, \dots, d_5 \in F.$$

Replace *s with any nonzero elements of F .

Example: Computing min rank in \mathbb{R}, F_2, F_3

$F = \mathbb{R}, F_3$:

$$A = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 2 & 2 & 1 \\ 1 & 1 & 2 & 2 & 1 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix}$$

$\text{rank } A = 2$, so $\text{mr}(\mathbb{R}, G) = 2$ and $\text{mr}(F_3, G) = 2$.

But in F_2 , $2 = 0$, so $A \notin S(F_2, G)$.

Example: Computing min rank in F_2

$F = F_2$:

Any $A \in S(F_2, G)$ has form $\begin{bmatrix} d_1 & 1 & 1 & 1 & 0 \\ 1 & d_2 & 1 & 1 & 0 \\ 1 & 1 & d_3 & 1 & 1 \\ 1 & 1 & 1 & d_4 & 1 \\ 0 & 0 & 1 & 1 & d_5 \end{bmatrix}$.

$A[145|235] = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & d_5 \end{bmatrix}$ has determinant 1 so $\text{rank } A \geq 3$.

Therefore $\text{mr}(F_2, G) \geq 3$.

Idea of Algorithm

To find the graphs characterizing $\{G \mid \text{mr}(F, G) \leq k\}$:

1. Construct all matrices A of rank $\leq k$ over F . Use the fact

$$A = U^t B U \iff \text{rank}(A) \leq k.$$

(B is $k \times k$, rank k ; U is $k \times n$.)

2. Return non-isomorphic graphs corresponding to matrices.

Problem: Too many matrices.

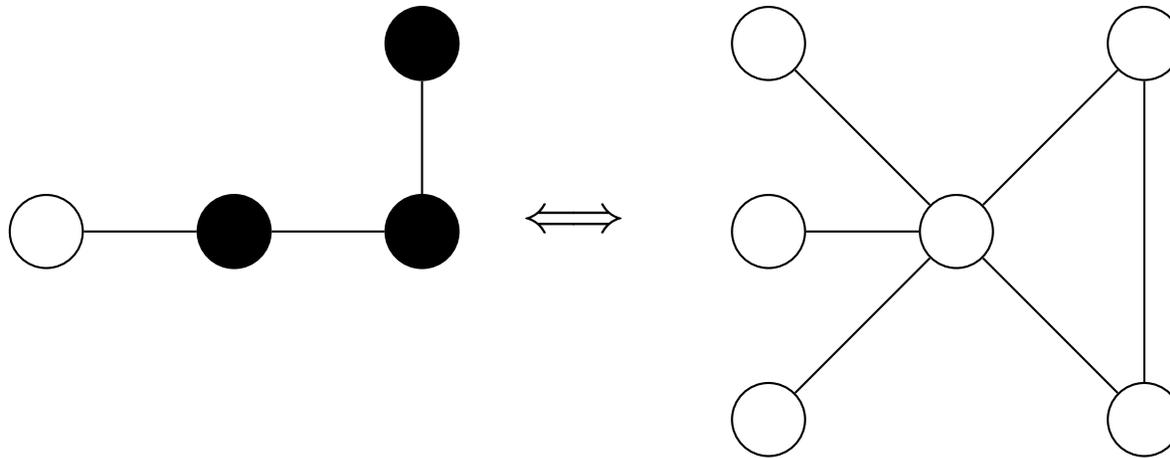
Solution: We only need the zero-nonzero patterns for the matrices. Be smarter by understanding $A = U^t B U$ better.

$$A = U^t B U$$

Feature/operation on U	Effect on graph corresponding to A
Column in U	vertex in graph
Column in U isotropic wrt B	zero entry for the vertex on diagonal
Column in U not isotropic wrt B	nonzero entry for the vertex on diagonal
Two columns orthogonal wrt B	no edge between corresponding vertices (zero matrix entry)
Two columns not orthogonal wrt B	edge between corresponding vertices (nonzero matrix entry)

$$A = U^t B U$$

Feature/operation on U	Effect on graph corresponding to A
Duplicate columns (isotropic)	independent set, vertices have same neighbors
Duplicate columns (non-isotropic)	clique, vertices have same neighbors
Columns multiples of each other	corresponding vertices have same neighbors (remember, only the zero-nonzero pattern is needed, and there are no zero divisors in F)
Interchanging two columns	relabel vertices



- black vertex \iff clique
- white vertex \iff independent set
- edge \iff all possible edges
- cliques or independent sets can be empty

Algorithm

To find the graphs characterizing $\{G \mid \text{mr}(F, G) \leq k\}$:

1. Find a maximal set of k -dimension vectors in F such that no vector is a multiple of any other. These are columns in U .
2. Construct all *interesting* matrices A of rank $\leq k$ over F . Use the fact

$$A = U^t B U \iff \text{rank}(A) \leq k.$$

(B is $k \times k$, rank k ; U is $k \times n$.)

3. Return non-isomorphic *marked* graphs corresponding to matrices.

Characterizing marked graphs for F_3

Rank	Vertices	Edges
2	5	5
2	5	4
3	14	54
4	41	525
4	41	528
5	122	4 860
6	365	44 100
6	365	44 109
7	1094	398 034

Finding forbidden subgraphs

Let S be the set of marked subgraphs of our characterizing graphs.

For each (normal) graph G

1. Construct the set T of possible marked graphs for G (can do this in exponential time).
2. If $S \cap T \neq \emptyset$, then G is a substitution graph of the characterizing graphs.
3. If $S \cap T = \emptyset$, then G is forbidden.

Open Questions

- Given a finite field F and positive integer k , what is a good upper bound for the number of vertices in minimal forbidden subgraphs?
- Is the bound 8 for $F = F_2$ and $k = 3$?
- Let G be any graph and let F be a finite field, $\text{char } F \neq 2$. Is $\text{mr}(\mathbb{R}, G) \leq \text{mr}(F, G)$? (true if $\text{mr}(\mathbb{R}, G) \leq 3$).